





Beyond the latency: New metrics for the real-time kernel

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In the beginning

In the begin a program was only a **logical sequence**, Then gosh said: we can't wait forever, we need to put **time** on this,

Since then we have two problems: The **logical correctness**, and the **timing correctness**.



In theory...

The systems defined as a set of tasks au Each task is a set of variables that defines its timing behavior, e.g.,

$$\mathcal{T}_{i} = \{P,C,D,B,J\}$$

Then, they try to define/develop a scheduler in such way that, for each task i in au:

the response time of $\tau_i < D_i$



For task level fixed priority scheduler:

 \forall task $i \in \tau$:

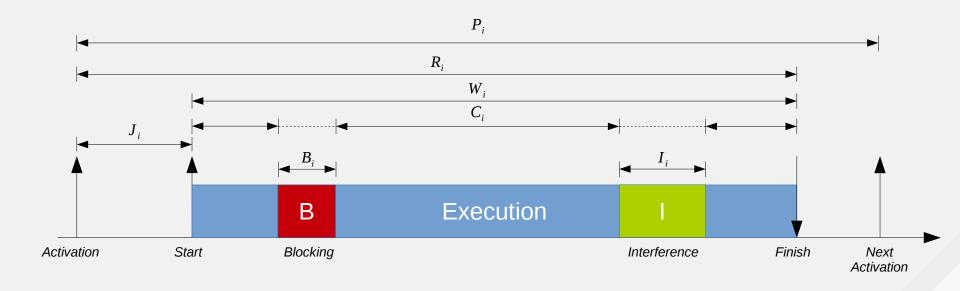
$$W_i = C_i + B_i + \sum_{j \in hp(i)} \left| \frac{W_i + J_j}{P_j} \right| C_j$$

$$R_i = W_i + J_i$$

is schedulable $\Leftrightarrow \forall task i \in \tau | R_i < D_i$



New metrics for the PREEMPT RT





PREEMPT_RT Timing correctness

- The preempt RT main metric is the latency
 - It is good, per carità...
- But it is very simplistic, if compared to response time.
- Latency is not even clearly defined
 - Kernel is seeing as a black box
 - There is no guarantee that the latency that took place now, will take place in the future (reproducibility/repeatability).
- It very hard, if not impossible, to give any guarantee in those numbers
 - We tried to use Extreme Value Analysis it does not fit in the method.



PREEMPT_RT Timing correctness

- User applications also depends on other characteristics of the kernel:
 - Locking
 - Dependence of other tasks
 - Interference of other tasks (and IRQs)

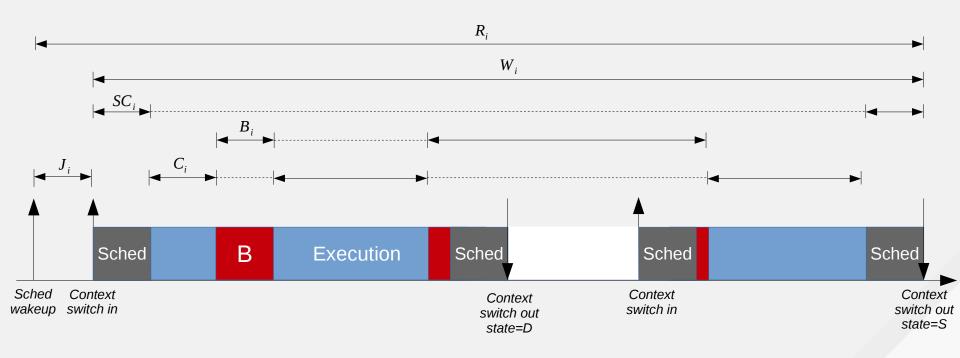


New metrics for the PREEMPT RT

- How can we improve the situation for Linux?
- What are tasks on Linux?
- What are the other metrics?
 - Execution time of task?
 - Blocking time? (SCHED_STATS)
 - Chain of locks that a task depends
 - Activation delay? (WAKEUP_DELAY)
 - Atomic context delay?
 - Dependency among tasks?



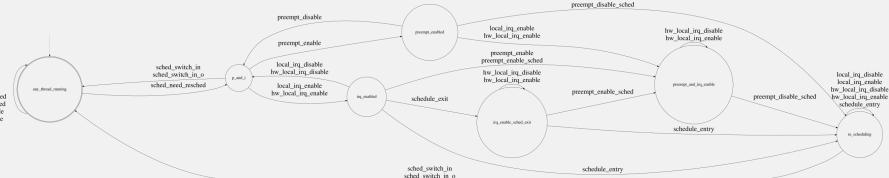
New metrics for the PREEMPT RT





What will I do, e.g., Composition of Latency

write abandon write_acquired write_blocked write lock mutex_abandon mutex_acquired mutex blocked mutex_lock read_abandon read_acquired read_blocked read_lock any_thread_running preempt disable sched preempt_enable_sched hw local irq disable hw_local_irq_enable local_irq_disable local_irq_enable preempt_disable preempt enable schedule entry schedule exit

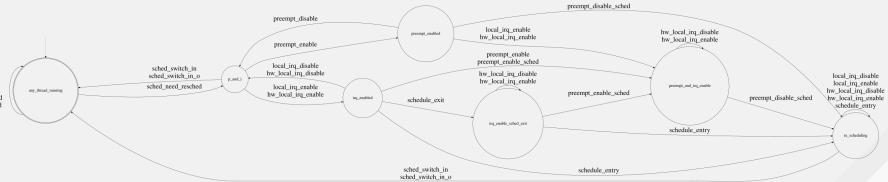




Rescheduling delay

- [need_resched...sched_return]
 - Case one: in the schedule

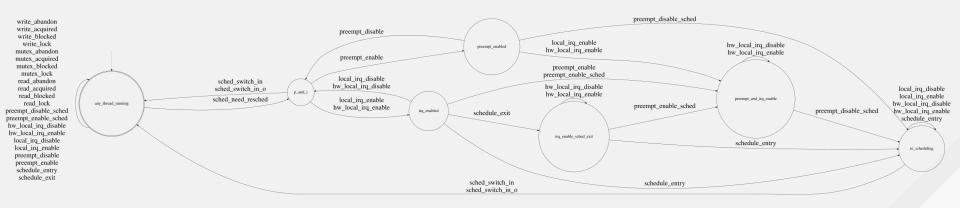
write abandon write_acquired write blocked write lock mutex_abandon mutex_acquired mutex_blocked mutex lock read_abandon read acquired read blocked read_lock preempt_disable_sched preempt_enable_sched hw_local_irq_disable hw_local_irq_enable local_irq_disable local_irq_enable preempt_disable preempt enable schedule_entry schedule exit





Rescheduling delay

- [need_resched...sched_return]
 - · Case two: calling the scheduler
 - Consider also that we have interference from interrupts





Thoughts?

- It is not reasonable doing this only in user-space
 - Too much data
- Should I do a trace-plugin?
- Use eBPF?
- Do something in kernel (lock stat like?)

